Math 6880/7875: Advanced Optimization Background and Review: Optimization

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Several topics are background for this course:

- (Numerical) linear algebra
- Probability/statistics
- "Basic" optimization knowledge

We'll spend some time briefly reviewing portions of these.

General optimization problem:

minimize f(x)subject to $x \in S$ $S := \{x \mid g_i(x) \leq 0, i \in [m]\}$

- $x = (x_1, \ldots, x_n)^T \in \mathbb{R}^n$ is the *optimization* or *design* variable
- $f: \mathbb{R}^n \to \mathbb{R}$ is the *objective* function
- $g_i : \mathbb{R}^n \to \mathbb{R}$, $i \in [m]$, are the *constraints*
- S is the *feasible* set
- We will always consider $n < \infty$, but we will occasionally allow $m \uparrow \infty$
- If m > 0, the problem is **constrained**; otherwise it is **unconstrained**

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$$(implicatly S = IR^{n})$$

General optimization problem:

minimize
$$f(x)$$

subject to $x \in S$
 $S := \{x \mid g_i(x) \leq 0, i \in [m]\}$

Solutions to optimization problems have their own taxonomy and properties:

- A point $x \in \mathbb{R}^n$ is feasible if $x \in S$.
- A point $x^* \in \mathbb{R}^n$ is a ("global") solution, optimum, or optimal point if $f(x^*) \leq f(x)$ for every $x \in S$.
- A point $x^* \in \mathbb{R}^n$ is a "local" solution, optimum, or optimal point if $\exists \epsilon > 0$ such that $f(x^*) \leq f(x)$ for all $x \in B_{\epsilon}(x^*) \cap S$, where

$$B_{\epsilon}(x^*) \cap S \coloneqq \left\{ x \in S \mid \|x - x^*\| < \epsilon \right\}.$$

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Solutions to optimization problems have their own taxonomy and properties:

- "Optimum"/"extremum" ↔ "maximum"/"minimum", as appropriate
- Maximization of f is minimization of -f
- Optimization problems can have zero, one, or many solutions. Which of these is true is rarely obvious.

Generally our goal is to find/compute an optimal solution. A local one could suffice.

~ equality constraints are double:
$$h(x)=0$$

 $g_1(x)=h(x)$
 $g_2(x)=-h(x)$

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Example One, unique solution



subject to $x \ge -1$

 $\min_{x \in \mathbb{R}} |x|$

Example

No solutions – infeasible

 $\min_{x \in \mathbb{R}} x^2$
subject to $|x| \leq -1$

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Example No solutions – unbounded

 $\max_{x \in \mathbb{R}} x^2$

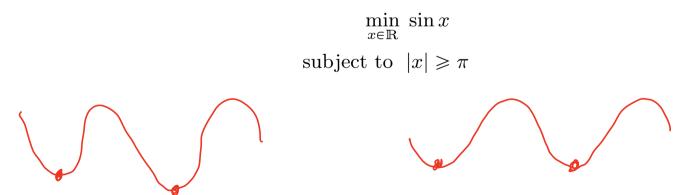
Example Many solutions

 $\min_{x \in \mathbb{R}} \sin x$
subject to $|x| \ge \pi$

Example No solutions – unbounded

 $\max_{x \in \mathbb{R}} x^2$





Ascertaining optimality

In some special cases, with some effort, one can conclude global optimality.

- Direct methods Analytically prove global optimality. (E.g., $f(x) = (x^2 1)^{10}$)
- Quadratic functions -f is quadratic with a positive-definite Hessian.
- Coercive functions Global optimality in some ball B, and show that that f outside B dominates f inside B.
- Globally convex functions Ensures that local minima are global minima.

Caveats:

- All the above are "easier" for unconstrained optimization, and become much more technical and difficult for constrained optimization.
- *Global* optimality requires some *global* knowledge of the objective and constraints.
- In high dimensions (*n* large), globally certifying any property of generic functions is hard.

The depressing fact of life: without relatively strong assumptions, local optimality is the best we know how to establish.

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There are a handful of **optimality conditions** that can be sufficient and/or necessary to determine local optimality.

- First-order optimality conditions conditions involving the gradients of f and/or g_i .
- Second-order optimality conditions conditions involving Hessians. Less computationally useful due to complexity/storage requirements.

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First-order local optimality: unconstrained optimization

Unconstrained optimization:

minimize f(x),

which implicitly allows $x \in \mathbb{R}^n$. (I.e., the feasible set is $S = \mathbb{R}^n$.)

The simplest first-order local optimality condition is a necessary one.

Theorem If $f \in C^1(\mathbb{R}^n)$, then x^* is a local minimum only if $\nabla f(x^*) = 0$.

Proof.

Fix $i \in [n]$. Let x_i be free, but fix $x_{\setminus i} = x^*_{\setminus i}$.

The resulting one-dimensional function f_i must have a local minimum at $x_i = x_i^*$, where its univariate derivative vanishes.

Repeat for every $i \Longrightarrow \nabla f(x^*) = 0$.

Notes:

- $\nabla f(x^*) = 0$ is not sufficient to conclude anything.
- ∇f(x*) = 0 is also a necessary condition for local minimization over S ⊂ ℝⁿ so long as x* ∈ int(S).

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Stationary points and definite matrices

Given $f : \mathbb{R}^n \to \mathbb{R}$ that is differentiable, a point x satisfying $\nabla f(x) = 0$ is a stationary point.

- Stationary points can be local/global minima.
- Stationary points can be local/global maxima.
- Stationary points can be saddle points (neither a maximum nor a minimum).

Many computational methods attempt to compute stationary points, even if we can't classify the result.

Stationary points are not necessarily easy to compute....

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Quadratic classification of matrices are needed for second-order conditions:

- A symmetric matrix A ∈ ℝ^{n×n} is positive definite if x^TAx > 0 for every x ∈ ℝⁿ\{0}.
 We write A > 0.
 (Equivalently, the inequality holds for all x with unit norm.)
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- Similar definitions for negative definite, and negative semi-definite. (A < 0, $A \le 0$, respectively)
- Matrices that are not positive/negative definite are **indefinite**.

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Second-order local optimality: unconstrained optimization

Unconstrained optimization:

minimize f(x),

An initial, necessary second-order condition:

Theorem Assume $f \in C^2(\mathbb{R}^n)$. If x^* is a local minimum, then $\nabla f^2(x^*) \ge 0$.

Proof sketch.

Take second-order Taylor expansion of f around x^* ,

$$f(x) \approx f(x^*) + \nabla f(x^*)^T (x - x^*) + \frac{1}{2} (x - x^*)^T \nabla^2 f(x^*) (x - x^*).$$

 x^* must be a stationary point for f, and the above holds for all x sufficiently close to x^* .

As before, this necessary condition holds if x^* is in the interior of a feasible set for a constrained optimization problem.

$$f(x) = x^{3}, x^{*} = 0, f''(x^{*}) = 0 \implies x^{*}$$
 is a local min

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Review

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Sufficient second-order optimality

Unconstrained optimization:

minimize f(x),

Theorem Assume $f \in C^2(\mathbb{R}^n)$. If x^* is a stationary point for f and $\bigvee_{x^*} f^2(x^*) > 0$, then x^* is a local minimum.

Proof sketch.

Another second-order Taylor expansion of f for x close to x^* :

$$f(x) \approx f(x^*) + \nabla f(x^*)^T (x - x^*) + \frac{1}{2} (x - x^*)^T \nabla^2 f(x^*) (x - x^*).$$

 x^* is a stationary point for f, and $(x - x^*)^T \nabla^2 f(x^*)(x - x^*) > 0$.

Optimality for constrained optimization

Constrained optimization:

minimize f(x)subject to $x \in S$ $S := \{x \mid g_i(x) \leq 0, i \in [m]\}$

One major complication with constrained vs. unconstrained optimization: local optima on the boundary of the feasible set must be handled with care.

Given a local optimum x^* , we divide [m] into *active* and *inactive* constraint sets: • $\mathcal{A}(x^*) = \{i \in [m] \mid g_i(x) = 0\}$ • $\mathcal{I}(x^*) = \{i \in [m] \mid g_i(x) < 0\}$

No feasible descent: at a local minimum x^* , we cannot find a direction for travel that simultaneously decreases f and all element of $g_{\mathcal{A}(x^*)}$.

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Constraint qualification

minimize f(x)subject to $x \in S$ $S \coloneqq \{x \mid g_i(x) \leq 0, i \in [m]\}$

To state (useful versions of) first-order optimality, we require an additional concept.

A local minimum x^* satisfies the linear independence constraint qualification (LICQ) condition if

 $\left\{\nabla g_i(x^*)\right\}_{i\in\mathcal{A}(x^*)},$

is a collection of linearly independent vectors.

The LICQ condition is used to strengthen necessary local optimality conditions.

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Constrained optimization: first-order optimality

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Theorem (Karush-Kuhn-Tucker)

Assume both f and g_i are $C^1(S)$ for every $i \in [m]$. Assume x^* is a local minimum of the above optimization problem that satisfies the LICQ condition. Then there exists a $\lambda \in \mathbb{R}^m$ such that (x^*, λ) satisfies,

$$\nabla f(x^*) + \sum_{i \in [m]} \lambda_i \nabla g_i(x^*) = 0 \qquad (Stationarity)$$
$$\lambda_i g_i(x^*) = 0, \ i \in [m] \qquad (Complementary Slackness)$$
$$g_i(x^*) \leq 0, \ i \in [m] \qquad (Primal feasibility)$$
$$\lambda_i \geq 0, \ i \in [m] \qquad (Dual feasibility)$$

The above are called the **KKT conditions**, and any point (x, λ) satisfying these conditions (even if x is not a local minimum) is a **KKT point**.

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Equality constraint (1):

$$h(x)=0 \longrightarrow g_{1}(x)=h \quad = 9, (x)=0 \quad = 9h(x)=0$$

$$g_{2}(x)=-h \quad g_{2}(x)=0 \quad \downarrow \downarrow$$

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$$g_{2}(x)=0$$

KKT conditions proof idea

One proof of the KKT conditions is a combination of three ideas/techniques:

• No feasible descent: We cannot find any direction $d \in \mathbb{R}^n$ such that all the following hold:

$$\nabla f(x^*)^T \boldsymbol{d} < 0$$

$$\nabla g_i(x^*)^T \boldsymbol{d} < 0, \qquad i \in \mathcal{A}(x^*)$$

• Theorems of the alternative: If there does not exist a d satisfying the above, then there must exist a $\lambda \in \mathbb{R}^{m+1}$ with positive components satisfying

$$\lambda_0 \nabla f(x^*) + \sum_{i \in \mathcal{A}(x^*)} \lambda_i \nabla g_i(x^*) = 0, \qquad \lambda_i = 0, \quad i \in \mathcal{I}(x^*)$$

In particular, the above exercises Gordan's Theorem of the alternative.

• Constraint qualification: Under LICQ, we can set $\lambda_0 = 1$ without loss.

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- are a lot more complicated than unconstrained optimization conditions
- extend to equality constraints (associated dual inequality constraints are always active)
- are also necessary with other (typically weaker) types of constraint qualification
- technically don't require constraint qualification (Fritz-John conditions), but this makes them less useful
- are really only explicitly used to analytically solve problems, but serve as the basis for some algorithms

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