

Integration/differentiation with polynomial approximations

MATH 6610 Lecture 27

November 18, 2020

Polynomial approximation

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And we've explored the accuracy of polynomial interpolation.

(Lebesgue)

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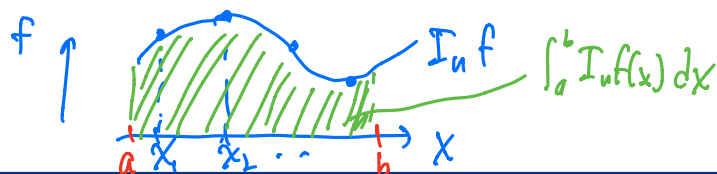
However, our main goal is utilizing polynomial methods for integration and differentiation.

The basic ideas:

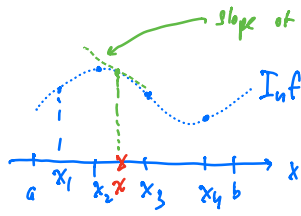
discrete approx. of an integral!

- Given function data at some points, we can construct a polynomial interpolant.
- (Interpolatory quadrature) Integration: we approximate the integral of the function as the integral of the polynomial interpolant.
- (Finite differences) Differentiation: we approximate the derivative of the function as the derivative of the polynomial interpolant.

$$\int_a^b f(x) dx \approx \int_a^b (I_n f)(x) dx \quad \leftarrow \text{can do exactly, } I_n f \in P_n.$$



Finite-difference:



$$(I_{inf} f)'(x) = \sum_{j=1}^4 w_j f(x_j)$$

Interpolatory quadrature

With x_1, \dots, x_n unique, fixed nodes on $[a, b]$, and f a given continuous function:
 Define $I_{n-1} : C([a, b]) \rightarrow P_{n-1}$ as the interpolation operator:

$$I_{n-1}f := \sum_{j=1}^n f(x_j) \ell_j(x),$$

$$= \sum_{j=1}^n c_j x^{j-1},$$

$$\ell_j(x) := \prod_{\substack{k=1, \dots, n \\ k \neq j}} \frac{x - x_k}{x_j - x_k}.$$

$$Vc = f. \quad (n \times n)$$

$$(V)_{j,k} = (x_j)^{k-1}$$

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$$= \sum_{j=1}^n c_j x^{j-1}, \quad Vc = f.$$

An *interpolatory quadrature* rule is a set of nodes $\{x_j\}$ and weights $\{w_j\}$ that results from using the integral of $I_{n-1}f$ to approximate the integral of f :

$$\int_a^b f(x)dx \approx \int_a^b [I_{n-1}f](x)dx =: \sum_{j=1}^n w_j f(x_j).$$

There are two ways to compute such nodes and weights, depending on the algorithmic strategy for I_{n-1} .

Interpolatory quadrature

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Example

Compute weights for the quadrature rule $\sum_{j=-1}^1 w_j f(j)$ that is accurate to as high a polynomial degree as possible. *(on the interval $[-1, 1]$).*

$$\int_{-1}^1 f(x) dx \approx w_{-1} f(-1) + w_0 f(0) + w_1 f(1).$$

2 ways:

1.) Enforce polynomial exactness

$$\int_{-1}^1 x^j dx = w_{-1} (-1)^j + w_0 (0)^j + w_1 (1)^j$$

for j as large as possible.

$$j=0: \int_{-1}^1 1 dx = w_{-1} (-1)^0 + w_0 \underbrace{(0)^0}_{\text{function } x^0 \equiv 1} + w_1 (1)^0$$

$$\begin{aligned} j=0: & 2 = w_{-1} + w_0 + w_1 \\ j=1: & 0 = -w_{-1} + w_1 \\ j=2: & \frac{2}{3} = w_{-1} + w_1 \end{aligned}$$

3 eqn's, 3 unknowns.

$$w_1 = w_{-1} \implies w_1 = \frac{1}{3} = w_{-1}$$

\Downarrow

$$w_0 = \frac{4}{3}$$

$$\int_{-1}^1 f(x) dx \approx \frac{1}{3} f(-1) + \frac{4}{3} f(0) + \frac{1}{3} f(1)$$

(exact $\forall f \in P_2$)

This procedure is equivalent to enforcing exact integration of an interpolant.

2.) Using Lagrange form.

$$\int_{-1}^1 f(x) dx \approx \int_{-1}^1 (I_{n-1} f)(x) dx = \sum_{j=-1}^1 w_j f(j).$$

(n=3)

$$(I_2 f)(x) = \sum_{j=-1}^1 f(j) \cdot l_j(x)$$

$$l_{-1}(x) = \frac{(x-0)(x-1)}{(-1-0)(-1-1)} = \frac{1}{2} x(x-1)$$

$$l_0(x) = \frac{(x+1)(x-1)}{(1)(-1)} = 1-x^2$$

$$l_1(x) = \frac{(x+1)x}{(2) \cdot 1} = \frac{1}{2} x(x+1)$$

$$\int_{-1}^1 (I_2 f)(x) dx = \sum_{j=-1}^1 f(j) \cdot \underbrace{\int_{-1}^1 l_j(x) dx}_{w_j}$$

$$\begin{aligned}w_{-1} &= \int_{-1}^1 l_{-1}(x) dx \\ &= \frac{1}{2} \left(\int_{-1}^1 x^2 dx - \int_{-1}^1 x dx \right) \\ &= \frac{1}{3}\end{aligned}$$

$$w_0 = \int_{-1}^1 (1-x^2) dx = 2 - \frac{2}{3} = \frac{4}{3}$$

$$w_1 = \int_{-1}^1 \frac{1}{2}(x^2+x) dx = \frac{1}{3}$$

$$\int_{-1}^1 f(x) dx \approx \int_{-1}^1 (\mathbb{I}_{n-1} f)(x) dx = \frac{1}{3} f(-1) + \frac{4}{3} f(0) + \frac{1}{3} f(1)$$

(same as before), (exact on P_2)

Interpolatory quadrature accuracy

The standard way to compute error estimates for polynomial quadrature: Taylor series arguments.

Typically one utilizes an error estimate for interpolation. For n -point interpolation on $[a, b]$, exact on P_{n-1} :

$$f(x) - [I_{n-1}f](x) = \frac{f^{(n)}(\xi)}{n!} \prod_{j=1}^n (x - x_j), \quad \xi = \xi(x) \in [a, b].$$

↑
from last time

$$\begin{aligned} \Rightarrow & \left| \int_a^b f(x) dx - \int_a^b (I_{n-1}f)(x) dx \right| \\ &= \left| \int_a^b \frac{f^{(n)}(\xi(x))}{n!} \prod_{j=1}^n (x - x_j) dx \right| \end{aligned}$$

note: $|x - x_j| \leq |b - a|$

$$\leq \int_a^b \frac{1}{n!} |f^{(n)}(\xi(x))| \cdot |b-a|^n dx$$

$$= \frac{(b-a)^n}{n!} \int_a^b |f^{(n)}(\xi(x))| dx$$

$$\leq \frac{(b-a)^{n+1}}{n!} \max_{x \in [a,b]} |f^{(n)}(x)|.$$

very small
if $b-a < 1$.

accuracy depends on smoothness

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Chaining this together with an integral:

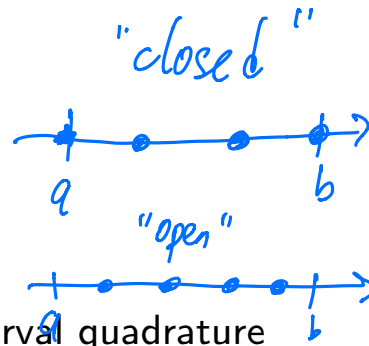
$$\left| \int_a^b f(x) dx - \int_a^b [I_{n-1}f](x) dx \right| \leq \frac{(b-a)^{n+1}}{n!} \max_{x \in [a, b]} |f^{(n)}(x)|.$$

Types of quadrature rules

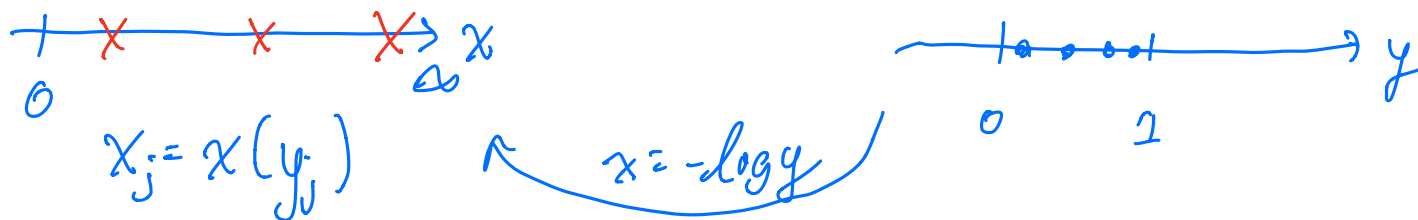
$$\int_a^b f(x) dx \approx \sum_{j=1}^n w_j f(x_j).$$

There are several special types of quadrature rules:

- Equidistant nodes: "Newton-Cotes" rules
 - ▶ "Closed": the endpoints are nodes (e.g., a, b)
 - ▶ "Open": all nodes are interior to the interval
- Change-of-variable tricks, e.g., used for infinite-interval quadrature



$$\int_0^{\infty} g(x) dx \approx \sum_{j=1}^n w_j f(x_j) \xrightarrow{y = e^{-x}} \int_0^1 f(y) dy \approx \sum_{j=1}^n w_j f(y_j)$$



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- Chebyshev/arcsine-like nodal arrangements: Clenshaw-Curtis quadrature, Fejér quadrature
- Hermite quadrature: uses derivatives as well as function values

$$\int_a^b f(x) dx \approx \sum_{j=1}^n w_j f(x_j) + \sum_{j=1}^n w'_j f'(x_j)$$

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- Chebyshev/arcsine-like nodal arrangements: Clenshaw-Curtis quadrature, Fejér quadrature
- Hermite quadrature: uses derivatives as well as function values
- Gaussian quadrature: Fixing n , has the maximum possible polynomial degree of exactness.

(Talk about Mon. after T-giving)

Finite difference formulas (approximations of derivatives) L27-S05

With x_1, \dots, x_n unique, fixed nodes on $[a, b]$, and f a given continuous function:
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A *finite difference formula* rule is a set of nodes $\{x_j\}$ and weights $\{w_j\}$ that results from using the derivative of $[I_{n-1}f](x)$ to approximate $f'(x)$.

Some minor differences from the quadrature case:

- The weights depend on the location x . (★)
- The weights can be defined by means other than interpolation. (Taylor Series)

Again, there are multiple ways to compute weights depending on how the interpolant is identified.

$$f'(x) \approx (I_{n-1}f)'(x) = \sum_{j=1}^n f(x_j) \overbrace{\ell_j'(x)}^{\text{weights}} \quad w_j = w_j(x)$$

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Example

Compute a 3-point finite-difference formula on the nodes $\{-1, 0, 1\}$ for differentiation at $x = 0$.

$$n=3, \quad \mathbb{I}_2 f(x) = \sum_{j=-1}^1 f(j) l_j(x)$$

$$l_{-1}(x) = \frac{1}{2} x(x-1)$$

$$l_0(x) = 1-x^2$$

$$l_1(x) = \frac{1}{2} x(x+1)$$

$$f'(0) \approx (\mathbb{I}_2 f)'(0) = \sum_{j=-1}^1 f(j) l_j'(0)$$

$$l_{-1}'(0) = (x - \frac{1}{2}) \Big|_{x=0} = -\frac{1}{2}$$

$$l_0'(0) = -2x \Big|_{x=0} = 0$$

$$l_1'(0) = (x + \frac{1}{2}) \Big|_{x=0} = \frac{1}{2}$$

$$f'(0) \approx \frac{1}{2} f(1) + 0 - \frac{1}{2} f(-1).$$

$$\text{nodes} = \{-1, 0, 1\}$$

$$\text{weights} = \{-\frac{1}{2}, 0, \frac{1}{2}\}.$$

Order of accuracy

Error estimates for finite difference formulas are computable via Taylor's Theorem. Recall that interpolation error on the interval $[a, b]$ scales like $(b - a)^{n+1}$.

Thus, finite difference formulas become more accurate as $(b - a) \rightarrow 0$.

Interest: as $(b-a) \rightarrow 0$, how fast does the finite-difference

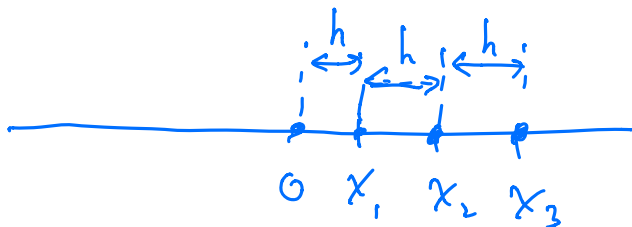
error go to zero?

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To standardize this, we typically assume equidistant nodes, with a spacing of $h > 0$. The *order of accuracy* of a finite difference formula is typically $\mathcal{O}(h^k)$ for some integer k .



k : convergence rate
 h^k : "truncation error"
 (and order of accuracy)

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Example

Compute a finite difference formula and order of accuracy (error estimate) for the three finite difference formulas:

$$(a) \quad f'(x) \approx w_1 f(x) + w_2 f(x + h) \quad (2\text{-pt formula})$$

$$f'(x) \approx w_1 f(x) + w_2 f(x - h)$$

$$(b) \quad f'(x) \approx w_1 f(x + h) + w_2 f(x - h) \quad (2\text{-pt formula})$$

(a) Could use Lagrange formula. Instead: Taylor Series.

$$f(x+h) = f(x) + h f'(x) + \frac{h^2}{2} f''(\xi) \quad \xi \in [x, x+h]$$

↑
Taylor

$\xi = \xi(x)$

$$w_1 f(x) + w_2 f(x+h)$$
$$= \underbrace{w_1 f(x) + w_2 f(x)}_{\text{want } = f'(x)} + \underbrace{h w_2 f'(x)}_{\text{want } = f'(x)} + \frac{h^2}{2} w_2 f''(\xi)$$

$$\left. \begin{aligned} (w_1 + w_2) f(x) = 0 &\Rightarrow w_1 + w_2 = 0 \\ h w_2 f'(x) = f'(x) &\Rightarrow h w_2 = 1 \end{aligned} \right\} \begin{aligned} w_2 &= \frac{1}{h} \\ w_1 &= -1/h \end{aligned}$$

$$w_1 f(x) + w_2 f(x+h) = f'(x) + \frac{h^2}{2} w_2 f''(\xi)$$

$$\stackrel{\text{for small } h}{=} f'(x) + O(h)$$

↑
truncation error
convergence rate: 1

(b) $f(x \pm h) = f(x) \pm h f'(x) + o(h^2)$

↑
Taylor

$$w_1 f(x+h) + w_2 f(x-h)$$

$$= f(x)(w_1 + w_2) + f'(x)(hw_1 - hw_2) + \cancel{(w_1 + w_2)} O(h^2)$$

$O(w_1 h^3) + O(w_2 h^3)$

want $f'(x)$

$$\left. \begin{array}{l} w_1 + w_2 = 0 \\ hw_1 - hw_2 = 1 \end{array} \right\} \begin{array}{l} w_2 = \frac{1}{2h} \\ w_1 = -\frac{1}{2h} \end{array}$$

$$w_1 f(x+h) + w_2 f(x-h) = f'(x) + O(w_1 h^2) + O(w_2 h^2)$$

$$= f'(x) + O(h)$$

rate of conv: 1

truncation error: $O(h)$

(this is wrong)

correct estimate:

$$f(x \pm h) = f(x) \pm hf'(x) + \frac{h^2}{2} f''(x) \pm \frac{h^3}{6} f'''(x) + O(h^4)$$

$$\begin{array}{ccc} \downarrow \frac{1}{2h} & & \downarrow -\frac{1}{2h} \\ w_1 f(x+h) + w_2 f(x-h) = f'(x) + \cancel{(w_1 + w_2)} \frac{h^2}{2} f''(x) \end{array}$$

$$+ (w_1 - w_2) \frac{h^3}{6} f'''(x)$$

$$+ O(h^4) w_1$$

$$= f'(x) + \underbrace{\frac{1}{h} - \frac{h^3}{6} f'''(x)}_{O(h^2)} + O(h^4) w_1$$

$$= f'(x) + O(h^2)$$

second-order accurate.

rate of conv: 2

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Odds and ends: higher-order derivatives can be approximated, function derivatives instead of function values can be used, etc.