

Integration/differentiation with polynomial approximations

MATH 6610 Lecture 27

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Polynomial approximation

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The basic ideas:

- Given function data at some points, we can construct a polynomial interpolant.
- (Interpolatory quadrature) Integration: we approximate the integral of the function as the integral of the polynomial interpolant.
- (Finite differences) Differentiation: we approximate the derivative of the function as the derivative of the polynomial interpolant.

Interpolatory quadrature

With x_1, \dots, x_n unique, fixed nodes on $[a, b]$, and f a given continuous function:
Define $I_{n-1} : C([a, b]) \rightarrow P_{n-1}$ as the interpolation operator:

$$\begin{aligned} I_{n-1}f &:= \sum_{j=1}^n f(x_j)\ell_j(x), & \ell_j(x) &:= \prod_{\substack{k=1, \dots, n \\ k \neq j}} \frac{x - x_k}{x_j - x_k}. \\ &= \sum_{j=1}^n c_j x^{j-1}, & Vc &= f. \end{aligned}$$

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$$= \sum_{j=1}^n c_j x^{j-1}, \quad Vc = f.$$

An *interpolatory quadrature* rule is a set of nodes $\{x_j\}$ and weights $\{w_j\}$ that results from using the integral of $I_{n-1}f$ to approximate the integral of f :

$$\int_a^b f(x)dx \approx \int_a^b [I_{n-1}f](x)dx =: \sum_{j=1}^n w_j f(x_j).$$

There are two ways to compute such nodes and weights, depending on the algorithmic strategy for I_{n-1} .

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Example

Compute weights for the quadrature rule $\sum_{j=-1}^1 w_j f(j)$ that is accurate to as high a polynomial degree as possible.

Interpolatory quadrature accuracy

The standard way to compute error estimates for polynomial quadrature: Taylor series arguments.

Typically one utilizes an error estimate for interpolation. For n -point interpolation on $[a, b]$, exact on P_{n-1} :

$$f(x) - [I_{n-1}f](x) = \frac{f^{(n)}(\xi)}{n!} \prod_{j=1}^n (x - x_j), \quad \xi = \xi(x) \in [a, b].$$

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Chaining this together with an integral:

$$\left| \int_a^b f(x) dx - \int_a^b [I_{n-1}f](x) dx \right| \leq \frac{(b-a)^{n+1}}{n!} \max_{x \in [a, b]} |f^{(n)}(x)|.$$

Types of quadrature rules

$$\int_a^b f(x)dx \approx \sum_{j=1}^n w_j f(x_j).$$

There are several special types of quadrature rules:

- Equidistant nodes: “Newton-Cotes” rules
 - ▶ “Closed”: the endpoints are nodes (e.g., a, b)
 - ▶ “Open”: all nodes are interior to the interval
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- Chebyshev/arcsine-like nodal arrangements: Clenshaw-Curtis quadrature, Fejér quadrature
- Hermite quadrature: uses derivatives as well as function values
- Gaussian quadrature: Fixing n , has the maximum possible polynomial degree of exactness.

Finite difference formulas

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A *finite difference formula* rule is a set of nodes $\{x_j\}$ and weights $\{w_j\}$ that results from using the derivative of $[I_{n-1}f](x)$ to approximate $f'(x)$.

Some minor differences from the quadrature case:

- The weights depend on the location x .
- The weights can be defined by means other than interpolation.

Again, there are multiple ways to compute weights depending on how the interpolant is identified.

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Example

Compute a 3-point finite-difference formula on the nodes $\{-1, 0, 1\}$ for differentiation at $x = 0$.

Order of accuracy

Error estimates for finite difference formulas are computable via Taylor's Theorem. Recall that interpolation error on the interval $[a, b]$ scales like $(b - a)^{n+1}$.

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Example

Compute a finite difference formula and order of accuracy (error estimate) for the three finite difference formulas:

$$f'(x) \approx w_1 f(x) + w_2 f(x + h)$$

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Odds and ends: higher-order derivatives can be approximated, function derivatives instead of function values can be used, etc.