

Polynomial approximation, II

MATH 6610 Lecture 26

November 13, 2020

Polynomial approximation

We've seen that polynomial interpolation is unisolvant. (in 1D)

Given a continuous f and distinct nodes x_1, \dots, x_{n+1} , then the Lagrange form of the interpolant is

$$p(x) = \sum_{j=1}^{n+1} f(x_j) \ell_j(x),$$

$$\ell_j(x) := \prod_{\substack{\ell=1, \dots, n+1 \\ \ell \neq j}} \frac{x - x_\ell}{x_j - x_\ell}.$$

(Lagrange form)

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We'll need some notation:

continuous functions on $[a, b]$

$$P_n := \text{span}\{1, x, \dots, x^n\}, \quad C([a, b]) := \{f : [a, b] \rightarrow \mathbb{R} \mid f \text{ is continuous on } [a, b]\},$$

and we metrize $C([a, b])$ with the norm,

$$\|f\| = \|f\|_\infty := \sup_{x \in [a, b]} |f(x)|.$$

Lebesgue functions and constants

To address accuracy, we investigate the *stability* of interpolation.

With x_1, \dots, x_{n+1} unique, fixed nodes, define $I_n : C([a, b]) \rightarrow P_n$ as the interpolation operator:

$$I_n f := \sum_{j=1}^{n+1} f(x_j) l_j(x). \quad (I_n \text{ is linear!})$$

(Note: I_n is a projector!) (i.e. $I_n I_n f = I_n f$)

Stability: $\|I_n\| = \sup_{f \neq 0} \frac{\|I_n f\|}{\|f\|} = \sup_{\|f\|=1} \|I_n f\|.$

↑
induced
norm

We can also look at stability of interpolation at a particular $x \in [a, b]$.

This notion of stability: $\sup_{\|f\|=1} |I_n f(x)|$

$$= \sup_{\|f\|=1} \left| \sum_{j=1}^{n+1} f(x_j) l_j(x) \right|$$

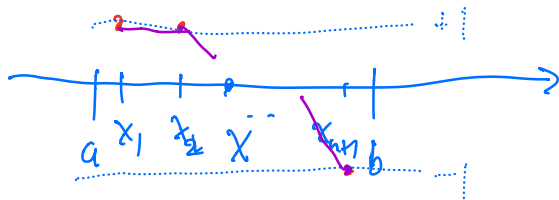
$$(\star) \leq \sup_{\|f\|=1} \sum_{j=1}^{n+1} |f(x_j)| \cdot |l_j(x)|$$

$$(\star\star) \leq \sum_{j=1}^{n+1} |l_j(x)|.$$

We can make this an equality:

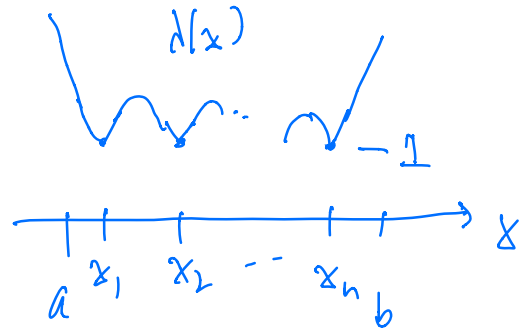
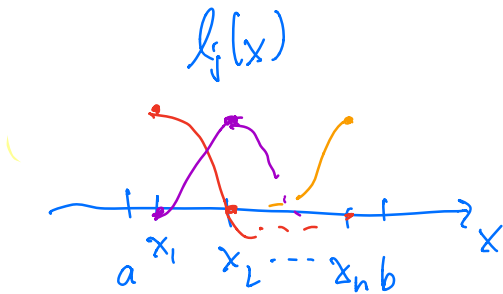
(\star): choose $f(x_j) = \text{sgn } l_j(x)$.

($\star\star$): pick f as pw-linear between x_j



$$S_0 := \sup_{\|f\|=1} |\mathcal{I}_n f(x)| = \sum_{j=0}^{n+1} |l_j(x)| := \lambda(x)$$

"Lebesgue function"



General interpolation stability:

$$\|\mathcal{I}_n\| = \sup_{\|f\|=1} \|\mathcal{I}_n f\| = \sup_{\|f\|=1} \|\lambda(x)\|$$

$$= \|\lambda(x)\| := \Lambda$$

"Lebesgue constant"

Lebesgue functions and constants

To address accuracy, we investigate the *stability* of interpolation.

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A computation shows that

$$\|I_n\|_{C([a,b]) \rightarrow P_n} = \Lambda := \|\lambda(x)\|_{\infty},$$

where

$$\lambda(x) := \sum_{j=1}^{n+1} |\ell_j(x)|$$

λ is called the Lebesgue function, and Λ the Lebesgue constant.

Stability gives us a way to investigate accuracy.

$$\begin{aligned}\|f - I_n f\| &= \|f - p + p - I_n f\| \quad (p \in P_n) \\ &\leq \|f - p\| + \|p - I_n f\|\end{aligned}$$

$$= \|f - p\| + \|I_n p - I_n f\| \quad (p = I_n p)$$

$$= \|f - p\| + \|I_n(p - f)\| \quad (I_n \text{ is linear})$$

$$\leq \|f - p\| + \|I_n\| \cdot \|p - f\| \quad (\text{true for all induced norms})$$

$$= (1 + \Lambda) \|f - p\|$$

$$\Rightarrow \|f - I_n f\| \leq \underbrace{(1 + \Lambda)}_{\substack{\text{depends only on nodes} \\ \text{independent of } f}} \cdot \underbrace{\inf_{p \in P_n} \|f - p\|}_{\substack{\text{best polynomial approx error.} \\ \text{independent of nodal choice.} \\ \text{(depends only on } f, n)}}.$$

depends only on nodes
independent of f .

best polynomial approx error.
independent of nodal choice.
(depends only on f, n).

Lebesgue's Lemma

With the operator norm of interpolation, here is a classical result quantifying the quality of polynomial interpolation:

Lemma (Lebesgue)

Let $f \in C([a, b])$, and assume that x_1, \dots, x_{n+1} are distinct nodes on $[a, b]$. Then

$$\|f - I_n f\|_\infty \leq (1 + \Lambda) \inf_{p \in P_n} \|f - p\|_\infty$$

(1 + Λ_n)

Thus, the Lebesgue constant governs the accuracy of polynomial interpolants.

Nodal sets

This raises the question: can we minimize Lebesgue constants?

Let Λ_n denote the Lebesgue constant for $n + 1$ points.

Λ_n and Λ_{n+1} need not have common nodes.

(We have a general triangular array of nodes:

$$\begin{array}{cccc}
 & & x_{1,1} & \leftarrow \Lambda_1 \\
 & & x_{1,2} & x_{2,2} \leftarrow \Lambda_2 \\
 & x_{1,3} & \dots & x_{2,3} \leftarrow \Lambda_3 \\
 / & & & \backslash \\
 & & & ; \\
 & & &)
 \end{array}$$

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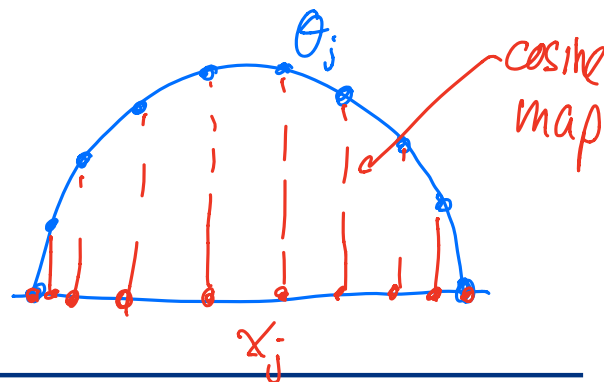
Λ_n and Λ_{n+1} need not have common nodes.

- For *any* triangular array of points, $\Lambda_n \xrightarrow{n \uparrow \infty} \infty$.
- For equidistant nodes, $\Lambda_n \sim 2^n$. (This is bad!)
- For "Chebyshev" nodes, $\Lambda_n \sim \log n$. (This is asymptotically optimal)

Chebyshev nodes x_j (say on $[-1, 1]$) are equidistant nodes under the cosine map:

$$x_j = \cos \theta_j, \quad \theta_j = \frac{2j-1}{2n+2} \pi, \quad j = 1, \dots, n+1.$$

(clusters nodes at boundary)



Another error estimate

There are alternative strategies to computing error estimates for interpolation.

With x_1, \dots, x_{n+1} the nodes, the *nodal polynomial* $\omega(x)$ is

$$\omega(x) := \prod_{j=1}^{n+1} (x - x_j).$$

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This can be used to construct an error estimate:

Theorem

Let x_1, \dots, x_{n+1} be distinct nodes on $[a, b]$. Then

$$|f(x) - [I_n f](x)| = \frac{|f^{(n+1)}(\xi)|}{(n+1)!} |\omega(x)|,$$

where $\xi = \xi(x)$ lies in $[a, b]$.