

## Polynomial approximation, II

MATH 6610 Lecture 26

November 13, 2020

## Polynomial approximation

We've seen that polynomial interpolation is unisolvant.

Given a continuous  $f$  and distinct nodes  $x_1, \dots, x_{n+1}$ , then the Lagrange form of the interpolant is

$$p(x) = \sum_{j=1}^{n+1} f(x_j) \ell_j(x), \quad \ell_j(x) := \prod_{\substack{\ell=1, \dots, n+1 \\ \ell \neq j}} \frac{x - x_\ell}{x_j - x_\ell}.$$

# Polynomial approximation

We've seen that polynomial interpolation is unisolvent.

Given a continuous  $f$  and distinct nodes  $x_1, \dots, x_{n+1}$ , then the Lagrange form of the interpolant is

$$p(x) = \sum_{j=1}^{n+1} f(x_j) \ell_j(x), \quad \ell_j(x) := \prod_{\substack{\ell=1, \dots, n+1 \\ \ell \neq j}} \frac{x - x_\ell}{x_j - x_\ell}.$$

Today: how accurate is this interpolant?

## Polynomial approximation

We've seen that polynomial interpolation is unisolvant.

Given a continuous  $f$  and distinct nodes  $x_1, \dots, x_{n+1}$ , then the Lagrange form of the interpolant is

$$p(x) = \sum_{j=1}^{n+1} f(x_j) \ell_j(x), \quad \ell_j(x) := \prod_{\substack{\ell=1, \dots, n+1 \\ \ell \neq j}} \frac{x - x_\ell}{x_j - x_\ell}.$$

Today: how accurate is this interpolant?

We'll need some notation:

$$P_n := \text{span}\{1, x, \dots, x^n\}, \quad C([a, b]) := \{f : [a, b] \rightarrow \mathbb{R} \mid f \text{ is continuous on } [a, b]\},$$

and we metrize  $C([a, b])$  with the norm,

$$\|f\| = \|f\|_\infty := \sup_{x \in [a, b]} |f(x)|.$$

## Lebesgue functions and constants

To address accuracy, we investigate the *stability* of interpolation.

With  $x_1, \dots, x_{n+1}$  unique, fixed nodes, define  $I_n : C([a, b]) \rightarrow P_n$  as the interpolation operator:

$$I_n f := \sum_{j=1}^{n+1} f(x_j) \ell_j(x).$$

(Note:  $I_n$  is a projector!)

## Lebesgue functions and constants

To address accuracy, we investigate the *stability* of interpolation.

With  $x_1, \dots, x_{n+1}$  unique, fixed nodes, define  $I_n : C([a, b]) \rightarrow P_n$  as the interpolation operator:

$$I_n f := \sum_{j=1}^{n+1} f(x_j) \ell_j(x).$$

(Note:  $I_n$  is a projector!)

A computation shows that

$$\|I_n\|_{C([a,b]) \rightarrow P_n} = \Lambda := \|\lambda(x)\|_{\infty},$$

where

$$\lambda(x) := \sum_{j=1}^{n+1} |\ell_j(x)|$$

$\lambda$  is called the Lebesgue function, and  $\Lambda$  the Lebesgue constant.

## Lebesgue's Lemma

With the operator norm of interpolation, here is a classical result quantifying the quality of polynomial interpolation:

### Lemma (Lebesgue)

Let  $f \in C([a, b])$ , and assume that  $x_1, \dots, x_{n+1}$  are distinct nodes on  $[a, b]$ . Then

$$\|f - I_n f\|_\infty \leq (1 + \Lambda) \inf_{p \in P_n} \|f - p\|_\infty$$

Thus, the Lebesgue constant governs the accuracy of polynomial interpolants.

## Nodal sets

This raises the question: can we minimize Lebesgue constants?

Let  $\Lambda_n$  denote the Lebesgue constant for  $n + 1$  points.

$\Lambda_n$  and  $\Lambda_{n+1}$  need not have common nodes.



## Nodal sets

This raises the question: can we minimize Lebesgue constants?

Let  $\Lambda_n$  denote the Lebesgue constant for  $n + 1$  points.

$\Lambda_n$  and  $\Lambda_{n+1}$  need not have common nodes.

- For any triangular array of points,  $\Lambda_n \xrightarrow{n \uparrow \infty} \infty$ .
- For equidistant nodes,  $\Lambda_n \sim 2^n$ . (This is bad!)
- For "Chebyshev" nodes,  $\Lambda_n \sim \log n$ . (This is asymptotically optimal)

Chebyshev nodes  $x_j$  (say on  $[-1, 1]$ ) are equidistant nodes under the cosine map:

$$x_j = \cos \theta_j, \quad \theta_j = \frac{2j-1}{2n+2} \pi, \quad j = 1, \dots, n+1.$$

## Another error estimate

There are alternative strategies to computing error estimates for interpolation.

With  $x_1, \dots, x_{n+1}$  the nodes, the *nodal polynomial*  $\omega(x)$  is

$$\omega(x) := \prod_{j=1}^{n+1} (x - x_j).$$

## Another error estimate

There are alternative strategies to computing error estimates for interpolation.

With  $x_1, \dots, x_{n+1}$  the nodes, the *nodal polynomial*  $\omega(x)$  is

$$\omega(x) := \prod_{j=1}^{n+1} (x - x_j).$$

This can be used to construct an error estimate:

### Theorem

Let  $x_1, \dots, x_{n+1}$  be distinct nodes on  $[a, b]$ . Then

$$|f(x) - [I_n f](x)| = \frac{|f^{(n+1)}(\xi)|}{(n+1)!} |\omega(x)|,$$

where  $\xi = \xi(x)$  lies in  $[a, b]$ .