L25-S00

# Polynomial approximation, I

MATH 6610 Lecture 25

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## Polynomial approximation

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## Polynomial approximation

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In the non-periodic case, perhaps the most popular alternative is polynomial approximation.

Like with Fourier Series, there is a "completeness" statement, ensuring that polynomials have sufficient approximation capacity for continuous functions.

#### Theorem (Weierstrass)

Let  $f : [a, b] \to \mathbb{R}$  be continuous with  $b - a < \infty$ . Then there exists a sequence of polynomials  $\{p_n\}_{n=0}^{\infty}$ , with deg  $p_n \leq n$ , such that

$$\lim_{n \uparrow \infty} \sup_{x \in [a,b]} |f(x) - p_n(x)| = 0.$$

This Weierstrass approximation result ensures that it's possible to construct accurate approximating polynomials.

### Polynomial interpolation

How do we construct polynomial approximations? Interpolation is the simplest strategy.

Given a continuous function  $f:[a,b] \to \mathbb{R}$  , we seek an interpolant from the space

$$P_{n} := \operatorname{span}\{1, x, \dots, x^{n}\}$$
on the points  $\{x_{1}, \dots, x_{n+1}\} \subset [a, b]$ .  
Constraints. Jegrees of freedom  
Generate  $p \in P_{n}$  s.t.  $p(x_{j}) = f(x'_{j}) \quad \forall j = 1, \dots, n+1$   

$$p(x) = \sum_{k=0}^{n} C_{KH} \times k$$

$$p(x_{i}) = f(x_{j}) \implies \sum_{k=0}^{n} c_{k+1} \chi_{j}^{k} = f(x_{i}^{*})$$

$$\downarrow$$

$$linear condition on \{c_{k}\}$$

constraints ?

#### Polynomial interpolation

How do we construct polynomial approximations? Interpolation is the simplest strategy.

Given a continuous function  $f:[a,b]\to \mathbb{R},$  we seek an interpolant from the space

$$P_n \coloneqq \operatorname{span}\{1, x, \dots, x^n\}$$

on the points  $\{x_1, \ldots x_{n+1}\} \subset [a, b]$ .

This results in a linear algebra problem: given  $p \in P_n$ , then

$$p(x) = \sum_{j=0}^{n} c_j x^j,$$

for some cofficients  $c_j$ . By enforcing equality on the interpolation grid:

$$p(x_j) = f(x_j),$$
  $j = 1, ..., n + 1,$ 

we obtain a linear system for the coefficients  $c_j$ :

$$Vc = f, \qquad (V)_{j,\ell} = x_j^{\ell-1}, \qquad | \leq (j,\ell) \leq (j,\ell)$$

and  $f_j = f(x_j)$ ,  $c_j = c_j$ .

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Interpolation existence and uniqueness

$$Vc = f,$$
  $(V)_{j,\ell} = x_j^{\ell-1},$ 

V is called a Vandermonde matrix.

Existence and uniqueness of the interpolation problem is equivalent to that of the linear system.

#### Theorem (Interpolation unisolvence)

If the nodes  $\{x_1, \ldots, x_{n+1}\} \subset [a, b]$  are all distinct, then given any continuous f, there is a unique degree-n polynomial that interpolates f on these nodes.

(not true as written in more than I dimension)

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#### Theorem (Interpolation unisolvence)

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There are (at least) two ways to prove this:

- Show that  $\det V \neq 0$ .
- Explicitly construct an interpolant (non-linear-algebraic).

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Proof: Uniqueness: assume 
$$p, q$$
 are both  
degree-  $n$  interpolants.  
 $p(X_6) = q(X_5)$ ,  $j = 1$ .  $n+1$   
 $X_j \neq X_k \forall j \neq k$   
 $\implies p(x) - g(x) \in P_n$   
 $p-q$  has  $n+1$  zeros  $@$   
 $X_1 - X_{n+1}$ .  
 $\implies Fund.$  theorem of algebra  
implies  $p-q \equiv 0$ ,  
 $E_{XiB-tenee}$   
(Strategy A) Show det  $V \neq 0$ ,  
 $lby$  induction).  
Lemma: If  $V_n$  is the nxn Vandermonde  
matrix:

$$V_{n} \equiv \begin{pmatrix} 1 & \chi_{i} & \chi_{i}^{2} - \chi_{i}^{n-1} \\ 1 & \chi_{L} & 1 & \ddots \\ \vdots & \vdots & 1 & \ddots \\ 1 & \chi_{n} & \chi_{n}^{2} - \chi_{n}^{n-1} \end{pmatrix}$$
Then  $= \det V_{n} \equiv \prod_{1 \leq j < k \leq n} (\chi_{k} - \chi_{j})$ .
  
Proof of lemma: (induction)
  
 $n \equiv j : V_{i} \equiv (1)$ ,  $\det V_{i} \equiv \prod_{1 \leq j < k \leq 1} (\chi_{k} - \chi_{j})$ 
  
 $= \prod$ 
  
 $\underline{n \equiv 2} = V_{2} = \begin{pmatrix} 1 & \chi_{i} \\ 1 & \chi_{2} \end{pmatrix}$ ,  $\det V_{2} \equiv \chi_{2} - \chi_{1}$ 
  
 $= \prod_{1 \leq j < k \leq 2} (\chi_{n} - \chi_{j})$ 
  
inductive step: assume  $\det V_{n} \equiv \prod (\chi_{k} - \chi_{j})$ 
  
 $\lim_{1 \leq j < k \leq 2} (1 - \chi_{1} - \chi_{n}^{n}) \lim_{1 \leq j < k \leq n} (\chi_{k} - \chi_{j})$ 

 $\det V_{at1} = \sum_{l=1}^{n+1} (-1) \chi_{n+1}^{l-1} M_{l}$   $Laplace exp
<math display="block"> along lass vow \qquad determinant mmor$ Me = Leteronivant of Vnm W/o row n+1 and column l. Note Me does not depend on Xn+1.  $\mathcal{M}_{\ell} = \mathcal{M}_{\ell}(\chi_{\ell}, \chi_{n})$ => bet Vin, as a function of Xing, is an element of Pn. As a function of Zny, then det Vny: • has roots @ X=X,...Xn. (replacing Xn+i by X; makes row mil of Vn+1 identical to row j => det Van=0).

• leading coefficient of det Vun is det Vn.  
det Vnti = 
$$\sum_{l=1}^{n+1} (-1)^{n+l-1} \chi_{nti}^{l-1} M_{l}^{l}$$
.  
=  $(-1)^{2n} \chi_{nti}^{n} M_{nti} + (lower order
terms)$   
=  $\chi_{nti}^{n} M_{nti} + \cdots$   
 $\int_{leterminist}^{n} determinist of Vnti with
row util, column n+1 removed
=  $\chi_{nti}^{n} (det V_n) + \cdots$   
 $det Vnti = C \cdot \prod_{l=1}^{n} (x_{nti} - x_l)$   
=  $(det V_n) \prod_{l=1}^{n} (x_{nti} - x_l)$   
=  $(\prod_{l \leq j \leq k \leq n} (x_k - x_j)) \prod_{l=1}^{n} (x_{nti} - x_l)$$ 

$$= \prod_{\substack{i \in j \in k \in N+1 \\ i \in j \in k \in N+1 \\ i \in j \in k \in N}} (X_{k} - X_{j}) \neq 0 \text{ if } X_{k} \neq X_{j}}$$
  
So: det  $V_{n} = \prod_{\substack{i \in j \in k \in N \\ i \in j \in k \in N}} (X_{k} - X_{0}) \neq 0 \text{ if } X_{k} \neq X_{j}}$   
 $\Rightarrow V_{c} = f \text{ has a (unique) solution.}$   
Strategy B (existence)  
Define  $l_{j}(x) = \prod_{\substack{k \in l = N \\ k \neq j}} (\frac{x - X_{k}}{(X_{j} - X_{k})} \in P_{n-1})$   
 $K_{e} \neq j$   
 $(j = [--N)$   
Note:  $l_{j}(X_{k}) = \begin{cases} 0 & j \neq k \\ 1 & j = k \end{cases}$   
 $l_{3} = 1$   
 $X_{1} = X_{2} = X_{3}$ 

. .

Then: 
$$\sum_{j=1}^{n} f(x_j) l_j(x) \in P_{n-1}$$
  
Satisfies  $\sum_{j=1}^{n} f(x_j) l_j(x_k)$   
 $= \sum_{j=1}^{n} f(x_j) \delta_{j,k} = f(x_k)$ 

The functions by (x) are called cardinal Lagrange functions.