# Department of Mathematics, University of Utah <br> Analysis of Numerical Methods I <br> MTH6610 - Section 001 - Fall 2017 

## Lecture notes

Monday October 16, 2017

These notes are not a substitute for class attendance. Their main purpose is to provide a lecture overview summarizing the topics covered.

Reading: Trefethen \& Bau III, Lecture 20
Reading: Isaacson \& Keller, Chapter 2.1
Let $A \in \mathbb{C}^{n \times n}$ be nonsingular. Given any $b \in \mathbb{C}^{n}$, there is a unique $x \in \mathbb{C}^{n}$ solving

$$
A x=b .
$$

The most elementary way of computing this solution is to solve the above system of $n$ equations and $n$ unknowns via Gaussian elimination. Practically, this entails forming the augmented matrix

$$
(A \vdots b),
$$

and subsequently performing row operations to bring the left $n \times n$ block into upper triangular form, followed by back-substitution to bring the left $n \times n$ block to the identity. The resulting vector in the last column is the solution $x$.
This procedure of performing row operations can be described without the need for the right-hand side vector $b$. In this way, we need not form an augmented matrix and instead just use $A$, but the price we pay is that we need to record the row operations that we use. This recording process can be codified by means of the $L U$ factorization of the matrix $A$. The matrix $A$ has columns

$$
A=\left(\begin{array}{llll}
a_{1} & a_{2} & \cdots & a_{n}
\end{array}\right), \quad \quad a_{j}=\left(\begin{array}{llll}
a_{j, 1} & a_{j, 2} & \ldots & a_{j, n}
\end{array}\right)^{T}
$$

For notational convenience, we define $a_{j}^{(1)} \equiv a_{j}$ for $j=1, \ldots, n$ and $A_{1} \equiv A$. The first step in Gaussian elimination would eliminate (zero-out) all entries in the column $a_{1}$ except for the first entry. If $r_{j}(A)$ represents the $j$ th row of $A$, this first elimination step performs

$$
\begin{aligned}
r_{1}(A) & \leftarrow r_{1}(A) \\
r_{2}(A) & \leftarrow r_{2}(A)-\frac{a_{1,2}}{a_{1,1}} r_{1}(A) \\
r_{3}(A) & \leftarrow r_{3}(A)-\frac{a_{1,3}}{a_{1,1}} r_{1}(A) \\
\vdots & \\
r_{n}(A) & \leftarrow r_{n}(A)-\frac{a_{1, n}}{a_{1,1}} r_{1}(A) .
\end{aligned}
$$

We can write this in matrix operations:

$$
\begin{aligned}
& A=\underbrace{\left(\begin{array}{llll}
\ell_{1} & e_{2} & \cdots & e_{n}
\end{array}\right)}_{L_{1}} \underbrace{\left(\begin{array}{llll}
a_{1}^{(2)} & a_{2}^{(2)} & \ldots & a_{n}^{(2)}
\end{array}\right)}_{A_{2}}, \\
& \ell_{1}=\left(\begin{array}{lllll}
1 & \frac{a_{1,2}}{a_{1,1}} & \frac{a_{1,3}}{a_{1,1}} & \cdots & \frac{a_{1, n}}{a_{1,1}}
\end{array}\right)^{T},
\end{aligned}
$$

where $a_{1}^{(2)}$ has zeros in entries $2, \ldots, n$.
We proceed by eliminating entries $3, \ldots, n$ from the vector $a_{2}^{(2)}$. This results in a similar formula:

$$
\begin{aligned}
& A_{2}=\underbrace{\left(\begin{array}{lllll}
e_{1} & \ell_{2} & e_{3} & \cdots & e_{n}
\end{array}\right)}_{L_{2}} \underbrace{\left(\begin{array}{lllll}
a_{1}^{(3)} & a_{2}^{(3)} & a_{3}^{(3)} & \ldots & a_{n}^{(3)}
\end{array}\right)}_{A_{3}}, \\
& \ell_{2}=\left(\begin{array}{llllll}
0 & 1 & \frac{a_{2,3}^{(2)}}{a_{2,2}^{(2)}} & \frac{a_{2,4}^{(2)}}{a_{2,2}^{(2)}} & \cdots & \frac{a_{2, n}^{(2)}}{a_{2,2}^{(2)}}
\end{array}\right)^{T},
\end{aligned}
$$

We can iterate this procedure: at step $k$, the first $k-1$ columns of $A_{k}$ are in upper triangular form. We then focus on the $k$ column of $A_{k}$, which is $a_{k}^{(k)}$. We form a matrix $L_{k}$ that eliminates entries $k+1, \ldots, n$ from $a_{k}^{(k)}$ by performing row operations. Putting this all together, we have

$$
A=L_{1} L_{2} \cdots L_{n-1} A_{n}
$$

Note that each $L_{k}$ is lower-triangular, with each column being a column of the identity, except for column $k$ which is the vector $\ell_{k}$. A computation with this structure of $L_{k}$ shows that

$$
L_{k} L_{k+1}=\left(\begin{array}{lllllllll}
e_{1} & e_{2} & \cdots & e_{k-1} & \ell_{k} & \ell_{k+1} & e_{k+2} & \cdots & e_{n}
\end{array}\right),
$$

which is also lower-triangular. Then making the definitions

$$
L:=\prod_{j=1}^{n-1} L_{j}, \quad U:=A_{n}
$$

we have the equality

$$
A=L U,
$$

with $L$ lower triangular and $U$ upper triangular. This is the $L U$ factorization of $A$. Note that, given $L$ and $U$, computing $A^{-1} b$ requires only forward- and back-substitution operations.
Note that not all invertible matrices have an $L U$ factorization: the steps above require us to divide by $a_{k, k}^{(k)}$ at step $k$. If any of these values vanish, then Gaussian elimination cannot be completed.

