# Topologically correct reconstruction of tortuous contour forests 

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## Outline

(1) Motivation and previous work

(2) Approach

(3) Results

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## Motivation - neuronal modeling



Fiala, Spacek, and Harris 2002


- Dendritic spine distribution and shape in the hippocampus correlates with certain types of brain disease.

A Neurologically normal (6 months).
B Mentally retarded (12 months).
C Alzheimer's (adult).
D Fragile $X$ syndrome (adult).

- Simulating the effect of spine modification requires accurate geometry and stable methods for modeling electric activity.
- This talk will focus on creating accurate meshed geometries suitable for electrophysiological finite-element simulations.


## Reconstruction

The high-level outline of a typical 3D reconstruction process is as follows:

- Begin with series of 2D Electron Microscopy (EM) images
- Generate contours around components of interest (axons, dendrites, etc in our case)
- Use some method to reconstruct 3D objects from 2D contours
- Combine 3D objects into a forest of structures


## Related work - single-component reconstruction


[6]

[3]

[1]

## Related work - multi-component reconstruction


[5]

[7]

[2]

- Recent work by Boissonnat and Memari [5] reconstructs single structures from non-parallel slices.
- Two approaches by Liu et al. [7] and Barequet and Vaxman [4] reconstruct from non-parallel slices and additionally reconstruct multiple components at the same time, avoiding inter-component intersections.
- Bajaj and Gillette [2] perform single component reconstruction using [1] and then remove intersections by removing contour overlaps in intermediate planes.


## Outline

## (1) Motivation and previous work

(2) Approach

## Single component reconstruction

With Bajaj's algorithm [1] we can produce nice single-component reconstructions. The reconstructions can be combined to produce a forest.


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## Single component reconstruction - properties


(2)

(3)

Properties that we will take advantage of:
(1) The reconstructed surface is a piecewise closed surface of polyhedra.
(2) Slicing the reconstructed surface on any of the original slices produces exactly the input contours.
(3) Any vertical (parallel to the $z$ axis) line segment between two adjacent slices intersects a single component exactly 0 or 1 times, or along exactly one line segment.

## Forest reconstruction

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## Inter-component intersections

The problem: intersections can occur between components.

Such intersections occur when data is

- highly anisotropic
- tightly packed
- tortuous



## Intersection removal

- Our algorithm removes intersections by adjusting only z-values of existing tiles (triangles)
(1) Can remove intersections without iterating through mid-slices
(2) Will not cause additional intersections
(3) Branching treated just like any other intersection
- Intersections often occur at branching points, but may occur where there is no branching. Our algorithm handles both intersections the same way.



## Penumbral contours

## Fact

All intersections occur in penumbral regions. A point's penumbral contour is the contour whose projection contains the projected point.


## Conflict points

A conflict point is a point of intersection. Somewhat more formally:

## Definition

Point $p^{g}$ is called a conflict point if there is some point $p^{y}$ such that the projections are equal $\left(p^{y \prime}=p^{g \prime}\right)$ and $p^{y}$ is closer to $p^{g}$ 's penumbral contour than $p^{g}$ is.


## Fact

Two components $C^{g}$ and $C^{y}$ intersect if and only if there is at least one conflict point on the surface of either component.

## Moving conflict points

Our algorithm will remove intersections without causing other intersections.

## Theorem

Moving any conflict point $p^{g}$ in the direction of its penumbral contour will not generate any additional conflict points among any pair of components.


Idea of proof: as a point moves toward its penumbral contour it won't enter any component because only two components can intersect in a given penumbra.

## Removing conflicts

We can resolve conflict points by moving them in the directions of their penumbral contours without worrying about causing additional intersections. Once all conflict points are resolved, all intersections are removed.


## Conflict removal algorithm

(1) Detect conflict points.
(2) Trace paths between conflict points along edges of yellow tile. We call these cut paths.
(3) Use original tiles and cut paths to induce new polygons.
(- Triangulate polygons and move conflict points along $z$-axis.


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(9) Triangulate polygons and move conflict points along z -axis.


## Separating by a given delta

$d=\frac{|\overline{\mathbf{A}} \times \overline{\mathbf{B}}|}{|\overline{\mathbf{B}}|}$
Substituting for $\overline{\mathbf{A}}$ and $\overline{\mathbf{B}}$ :

$$
\begin{aligned}
d^{2}= & \left(\left(A_{y}\left(B_{z}-\epsilon\right)-\left(A_{z}+\epsilon\right) B_{y}\right)^{2}\right. \\
+ & \left(\left(A_{z}+\epsilon\right) B_{x}-A_{x}\left(B_{z}-\epsilon\right)\right)^{2} \\
+ & \left.\left(A_{x} B_{y}-A_{y} B_{x}\right)^{2}\right) /\left(B_{x}^{2}+B_{y}^{2}+\left(B_{z}+\epsilon\right)^{2}\right)
\end{aligned}
$$



After collecting $\epsilon$ :

$$
\begin{aligned}
0= & \epsilon^{2}\left(\left(A_{y}+B_{y}\right)^{2}+\left(A_{x}+B_{x}\right)^{2}-d^{2}\right) \\
+ & \epsilon(2)\left(\left(A_{x}+B_{x}\right)\left(A_{z} B_{x}-A_{x} B_{z}\right)\right. \\
& \left.-\left(A_{y}+B_{y}\right)\left(A_{y} B_{z}-A_{z} B_{y}\right)-d^{2} A_{z}\right) \\
+ & \left(A_{y} B_{z}-A_{z} B_{y}\right)^{2}+\left(A_{z} B_{x}-A_{x} B_{z}\right)^{2} \\
& +\left(A_{x} B_{y}-A_{y} B_{x}\right)^{2}-d^{2}\left(B_{x}^{2}+B_{y}^{2}+B_{z}^{2}\right)
\end{aligned}
$$



## Separating by a given delta



## Theorem

$\epsilon<\left|p^{g}-\mathscr{Z}\left(p^{g}\right)\right|$ and $\epsilon<\left|p^{y}-\mathscr{Z}\left(p^{y}\right)\right|$
Idea of proof: as points approach original contours, which are separated by $d$, the chords will be separated by at least $d$ in the limit.

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## Illustrative examples



## Illustrative examples



## Results



## Results



## Results



## Conclusions and notes

- Algorithm is $O\left(n^{2}\right)$ where $n$ is the number of tiles.
- Average case is closer to $n \log n$ complexity of sweep line algorithm as large majority of 2D intersections are not conflict points.
- Original contours remain unchanged - only makes changes in interpolated data between slices
- Topologically correct and water tight
- Generates large number of extra triangles in intersecting regions



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