

# CS 6170: Computational Topology, Spring 2019

## Lecture 28

### Topological Data Analysis for Data Scientists

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# Cohomology

## Homology: elementary chain

- Simplicial complex  $K$ , e.g. a triangulation of an annulus.
- Vertices, edges and triangles:  $K^0 = \{v_i\}$ ,  $K^1 = \{e_i\}$  and  $K^2 = \{\Delta_i\}$ .
- By abuse of notation, each 0-, 1- and 2-simplex in  $K$  corresponds to an *elementary chain* of the same dimension, denoted as  $v_i$ ,  $e_i$  and  $\Delta_i$ .
- 0-, 1- and 2-chains: formal sums of 0-, 1- and 2-simplexes (elementary chains) with integer coefficients  $\mathbb{Z}_2 = \{0, 1\}$ :

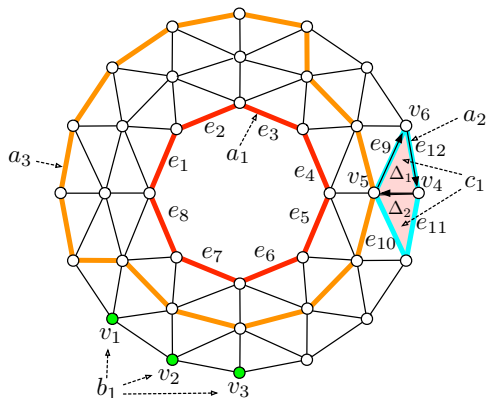
$$C_0 = C_0(K; \mathbb{Z}_2) = \{b = \sum g_i v_i \mid g_i \in \mathbb{Z}_2\},$$

$$C_1 = C_1(K; \mathbb{Z}_2) = \{a = \sum g_i e_i \mid g_i \in \mathbb{Z}_2\},$$

$$C_2 = C_2(K; \mathbb{Z}_2) = \{c = \sum g_i \Delta_i \mid g_i \in \mathbb{Z}_2\}.$$

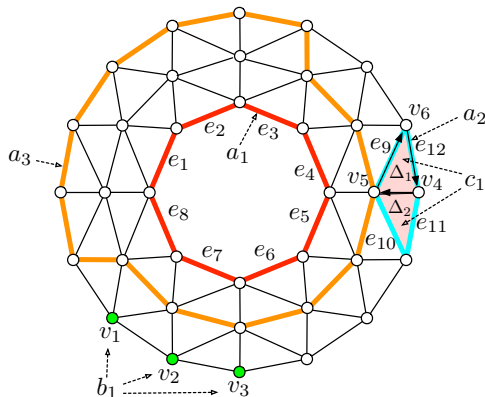
- Two elements  $a, a' \in C_1$  are homologous iff  $a - a' = \partial(c)$ , for some 2-chain  $c$ , denoted as  $a \sim a'$ . In this case  $[a] = [a']$ .

# Homology



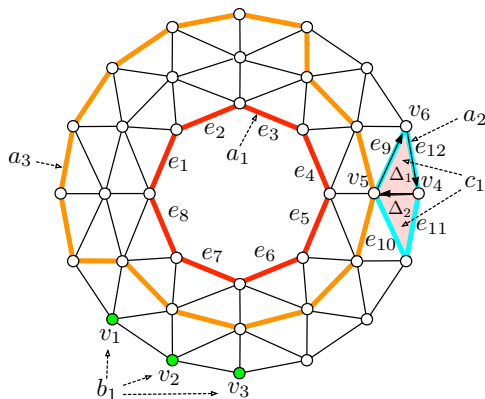
- 0-chain:  $b_1 = v_1 + v_2 + v_3$  (solid green),
- 1-chain:  $a_1 = e_1 + e_2 + e_3 + e_4 + e_5 + e_6 + e_7 + e_8$  (bold red),
- 2-chain:  $c_1 = \Delta_1 + \Delta_2$  (bold pink).
- Here  $a_1$  (bold red)  $\sim a_3$  (bold orange): they bound the “same” tunnel.

# Homology



- *Boundary maps:*  $\partial_2 : C_2 \rightarrow C_1$  and  $\partial_1 : C_1 \rightarrow C_0$
- *1-dimensional homology groups:*  $H_1 = Z_1/B_1 = \ker \partial_1 / \text{im} \partial_2$

# Homology



- $a_1$  (bold red) is a 1-cycle since  $\partial(a_1) = 0$ .
- $a_2 = e_9 + e_{10} + e_{11} + e_{12}$  (bold cyan) is a 1-boundary since it is the boundary of the 2-chain  $c_1$ .
- $a_1$  is a 1-cycle, but not a 1-boundary, which makes  $[a_1]$  a non-trivial element of  $H_1$ .

## Cohomology: elementary cochain

- Simplicial complex  $K$ , e.g. a triangulation of an annulus.
- Homology groups are “dual” to homology groups
- Vertices, edges and triangles:  $K^0 = \{v_i\}$ ,  $K^1 = \{e_i\}$  and  $K^2 = \{\Delta_i\}$ .
- By abuse of notation, each 0-, 1- and 2-simplex in  $K$  corresponds to an *elementary cochain* of the same dimension, denoted as  $v_i^*$ ,  $e_i^*$  and  $\Delta_i^*$ .
- 1-simplex  $e$  has a corresponding elementary 1-cochain  $e^*$ , which is a function on 1-chain whose value is 1 on  $e$  and 0 on all other edges.
- $e^* : C_1 \rightarrow \mathbb{Z}_2$ , where  $e^*(e) = 1$  and  $e^*(e') = 0$  for all  $e' \in K^1, e' \neq e$ .
- Similarly, we have elementary 0-cochains,  $v^*$  associated with the 0-simplices  $v$ ; and elementary 2-cochains  $\Delta^*$  associated with the 2-simplices  $\Delta$ .
- 0-, 1- and 2-cochains can be considered as sums of elementary cochains.

- 0-, 1- and 2-cochains: functions on 0-, 1- and 2-chain groups.

$$C^0 = C^0(K; \mathbb{Z}_2) = \{\beta : C_0 \rightarrow \mathbb{Z}_2, \beta = \sum g_i v_i^* \mid g_i \in \mathbb{Z}_2\},$$

$$C^1 = C^1(K; \mathbb{Z}_2) = \{\alpha : C_1 \rightarrow \mathbb{Z}_2, \alpha = \sum g_i e_i^* \mid g_i \in \mathbb{Z}_2\},$$

$$C^2 = C^2(K; \mathbb{Z}_2) = \{\gamma : C_2 \rightarrow \mathbb{Z}_2, \gamma = \sum g_i \Delta_i^* \mid g_i \in \mathbb{Z}_2\}.$$



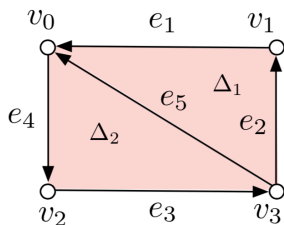
## Cohomology: coboundary map

- *Coboundary maps* are dual to the boundary maps,  $\delta_0 : C^0 \rightarrow C^1$ ,  $\delta_1 : C^1 \rightarrow C^2$
- Let  $\beta \in C^0$ ,  $\alpha \in C^1$ , we have:

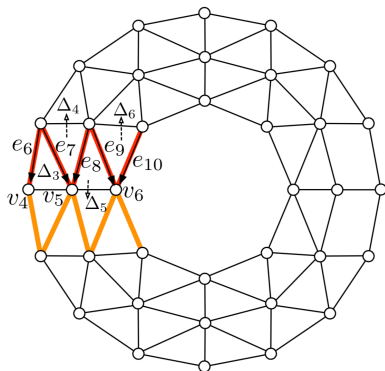
$$\begin{aligned}(\delta_0\beta)([v_0, v_1]) &= \beta(\partial_1([v_0, v_1]) = \beta(v_1) + \beta(v_0), \\(\delta_1\alpha)([v_0, v_1, v_2]) &= \alpha(\partial_2([v_0, v_1, v_2])) \\ &= \alpha([v_1, v_2]) + \alpha([v_0, v_2]) + \alpha([v_0, v_1]).\end{aligned}$$

- If  $\alpha = \sum g_i e_i^*$ , then  $\delta(\alpha) = \sum g_i (\delta e_i^*)$ .
- To compute  $\delta e^*$  for each oriented simplex  $e$ , we have  $\delta e^* = \sum \Delta_j^*$ , where the summation extends over all  $\Delta_j$  having  $e$  as a face.

- For a cochain  $\alpha \in C^1$ , we call  $\alpha$  a 1-cocycle if  $\delta_1(\alpha) = 0$ .
- We call  $\alpha$  a 1-coboundary if there exists a cochain  $\beta \in C^0$  such that  $\delta_0(\beta) = \alpha$ .
- It is easy to verify that  $\delta \circ \delta = 0$ .
- 1-coboundaries are always 1-cocycles, we have  $\text{im}(\delta_0) \subseteq \ker(\delta_1)$ .
- We define the 1-cohomology of  $K$  to be the quotient group,  $H^1 = Z^1/B^1 = \ker(\delta_1)/\text{im}(\delta_0)$ .
- Two 1-cocycles  $\alpha$  and  $\alpha'$  are cohomologous if  $\alpha + \alpha'$  is a coboundary.

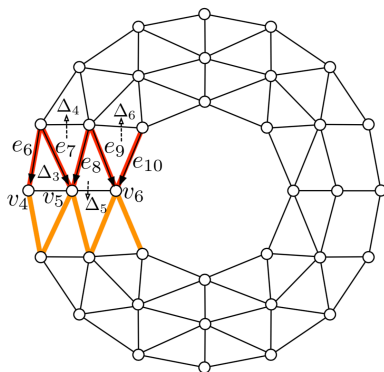


- $e_5^* : C_1 \rightarrow \mathbb{Z}$  has value 1 on  $e_5$  and 0 on other edges. Ignore orientations (for now).
- $\delta e_5^*$  has values 1 on  $\Delta_1$  and  $\Delta_2$ , because  $e_5$  appears in  $\partial\Delta_2$  and  $\partial\Delta_1$ .
- $\delta e_5^* = \Delta_2^* + \Delta_1^*$ .
- $\delta v_1^* = e_2^* + e_1^*$ ,  $\delta v_3^* = e_3^* + e_2^* + e_5^*$ .
- 1-cochain  $\alpha = e_1^* + e_5^* + e_3^*$  is a 1-cocycle since  $\delta(\alpha) = \delta(e_1^*) + \delta(e_5^*) + \delta(e_3^*) = (\Delta_1^*) + (\Delta_2^* + \Delta_1^*) + (\Delta_2^*) = 0$ .
- $\alpha$  is also a 1-coboundary since  $\alpha = \delta(v_1^* + v_3^*)$ .



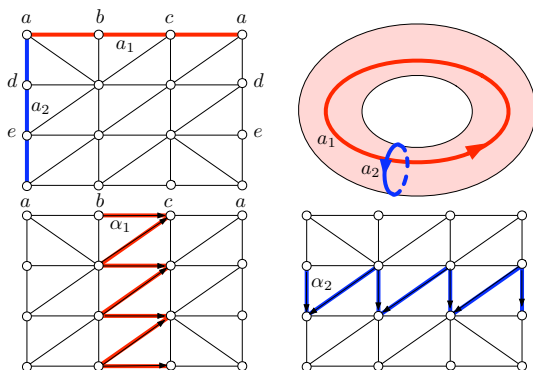
- 1-chain  $\alpha_1 = e_6^* + e_7^* + e_8^* + e_9^* + e_{10}^*$  is a 1-cocycle

$$\begin{aligned} \delta(\alpha_1) &= \delta(e_6^*) + \dots + \delta(e_{10}^*) \\ &= \Delta_3^* + (\Delta_4^* + \Delta_3^*) + (\Delta_5^* + \Delta_4^*) + (\Delta_6^* + \Delta_5^*) + \Delta_6^* = 0. \end{aligned}$$



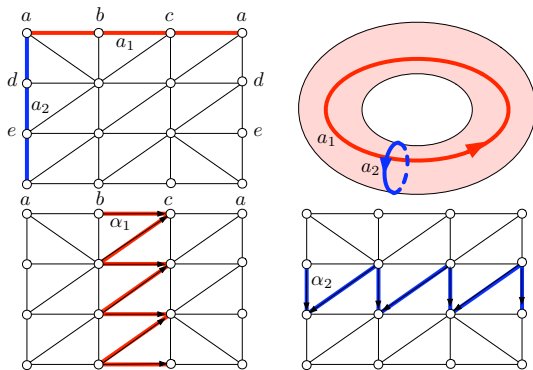
- 1-chain  $\alpha_1 = e_6^* + e_7^* + e_8^* + e_9^* + e_{10}^*$  is not a 1-coboundary.
- $[\alpha_1] \in H^1$ , and  $\alpha_1$  can be used as the representative of the 1-dimensional cohomology class.
- $\alpha_1$  (bold red) is cohomologous to  $\alpha_2$  (bold orange), as we can check  $\alpha_1 + \alpha_2 = \delta(v_4^* + v_5^* + v_6^*)$ .

# Homology groups of a torus



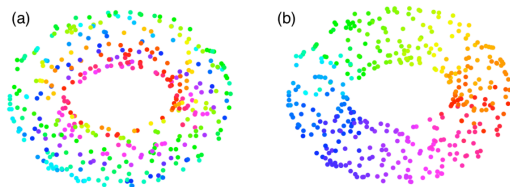
- $H_1$  is generated by the 1-chains  $a_1$  (red) and  $a_2$  (blue).
- $a_1 = [a, b] + [b, c] + [c, a]$  and  $a_2 = [a, d] + [d, e] + [e, a]$ .
- $a_1$  and  $a_2$  are 1-cycles, as  $\partial(a_1) = \partial(a_2) = 0$ .
- $a_1$  and  $a_2$  are not 1-boundaries.
- In addition,  $a_1$  and  $a_2$  are not homologous.

# Cohomology groups of a torus



- $H^1$  is generated by the 1-cochains  $\alpha_1$  (red) and  $\alpha_2$  (blue).  $\alpha_1$  and  $\alpha_2$  are 1-cocycles, not 1-coboundaries, and are not cohomologous.
- The duality between cohomology and homology generators is counter-intuitive.
- Here,  $\alpha_1 \in H^1$  (bold red) is dual to  $a_1 \in H_1$  (bold red), while  $\alpha_2 \in H^1$  (bold blue) is dual to  $a_2 \in H_1$  (bold blue).

# Cohomology based parametrization



de Silva et al. (2009)

- Using a principle from homotopy theory: relates circular coordinates with cohomology.
- Let  $[X, \mathbb{S}^1]$  be the set of equivalence classes of continuous maps from space  $X$  to  $\mathbb{S}^1$  under the homotopy relation.
- For topological spaces with the homotopy type of a cell complex, there is an isomorphism  $H^1(X; \mathbb{Z}) \cong [X, \mathbb{S}^1]$
- This implies that if  $X$  has a non-trivial 1-dimensional cohomology class  $[\alpha] \in H^1(X; \mathbb{Z})$ , we can construct a continuous function  $\theta : X \rightarrow \mathbb{S}^1$  from a representative  $\alpha$ .



## Future Directions and Discussions

What is the most valuable tool  
you have learnt in this class?

What is on your wish list?

What do you think are the future directions for topological data analysis?

- Topology in visualization: vector field topology, tensor field topology
- TDA for biomedicine: high-dimensional data analysis, mapper
- TDA for materials science, astronomy, music, signal processing
- Multi-parameter persistent homology
- Discrete Morse theory, discrete Stratified Morse theory

## Future directions

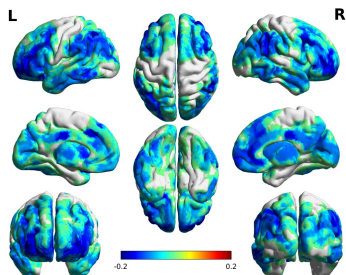
- Better understanding of multi-parameter persistent homology
- Scalable computation
- Machine learning
- New visualization tools, uncertainty visualization
- New theory leads to new algorithms and applications!

# Topological Data Analysis of Functional MRI Connectivity in Time and Space Domains

Anderson et al. (2018), Best Paper at MICCAI CNI 2018.

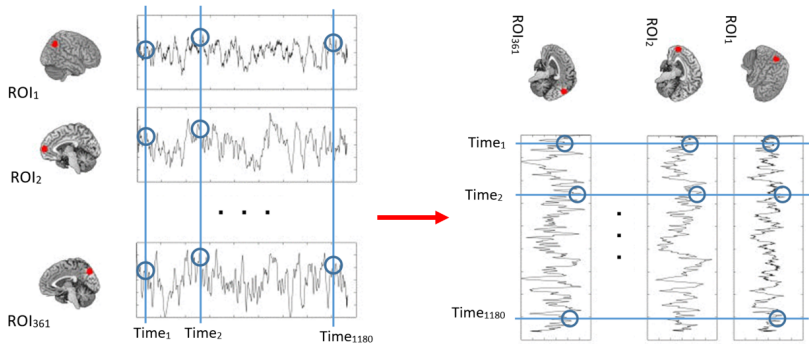
# Key findings

- Functional connectivity in time and space domains produced complementary information about brain function.
- Results from topological data analysis is significantly correlated with cognitive performance, especially fluid intelligence.
- Time and Space topology are both correlated with fluid intelligence, but the spatial domain picks up additional behavior and allows us to localize brain regions.



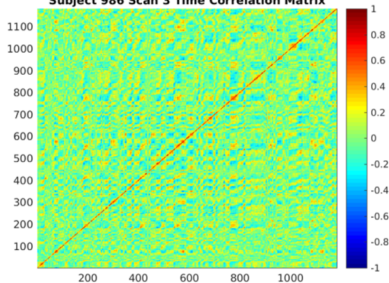


# Flipping the time series

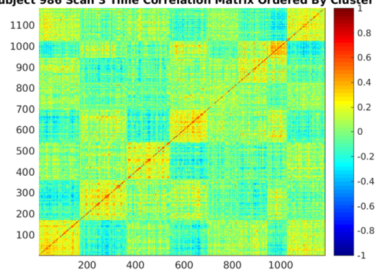


# Time domain connectivity matrix

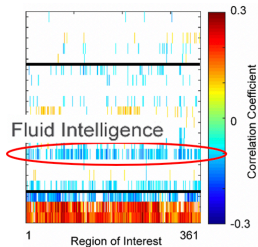
Subject 986 Scan 3 Time Correlation Matrix



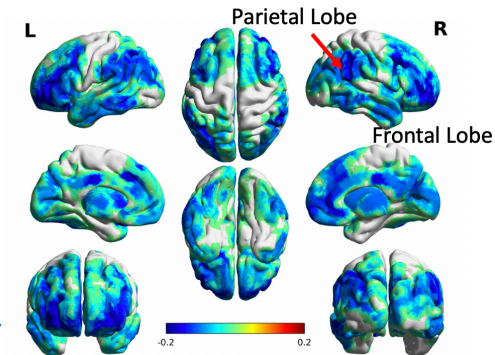
Subject 986 Scan 3 Time Correlation Matrix Ordered By Cluster



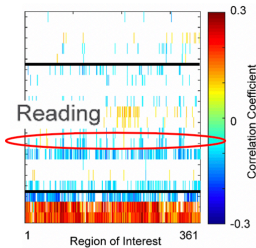
# TDA: Significant correlation with fluid intelligence



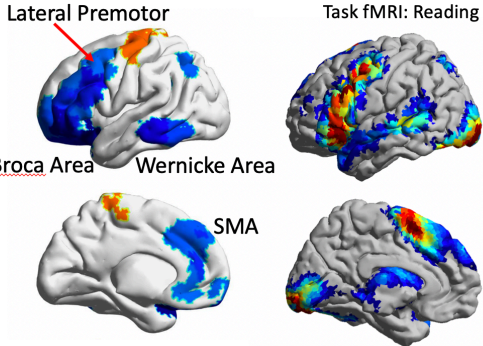
*Correlation matches  
fluid intelligence areas*



# TDA: Significant correlation with reading comprehension



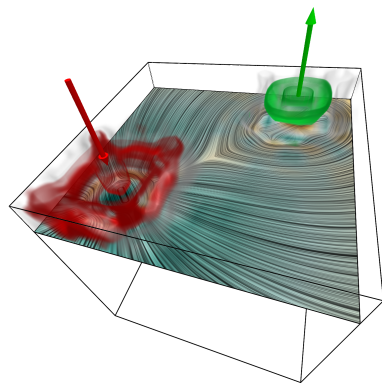
*Correlation matches  
key language areas*



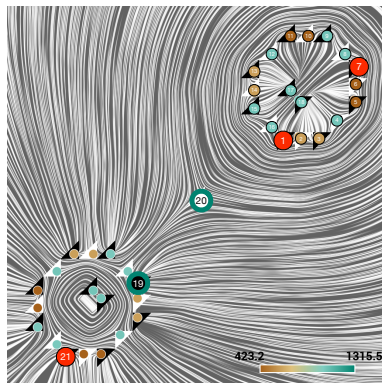
# Extract Robust Features From Stress Tensor Fields

Wang and Hotz (2017); Jankowai et al. (2018)

# Stress tensor field: two point load



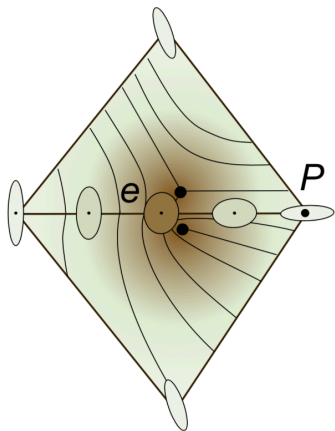
(a)



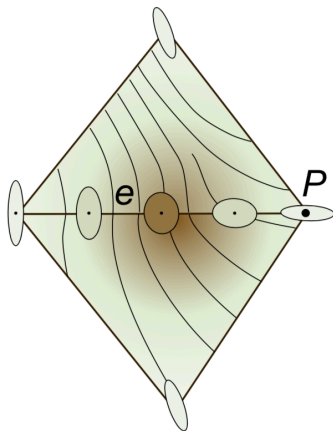
(b)

**Figure:** Visualization of the robustness of a slice of a stress tensor field. (a) A single slice of the data embedded in a 3D context visualized using volume rendering. (b) A textured slice with degenerate cells. White triangles represent trisectors and black ones represent wedges. Degenerate points in each cell are visualized with a brown-to-turquoise colormap. Degenerate points with infinite robustness are in red.

# Robustness of degenerate points

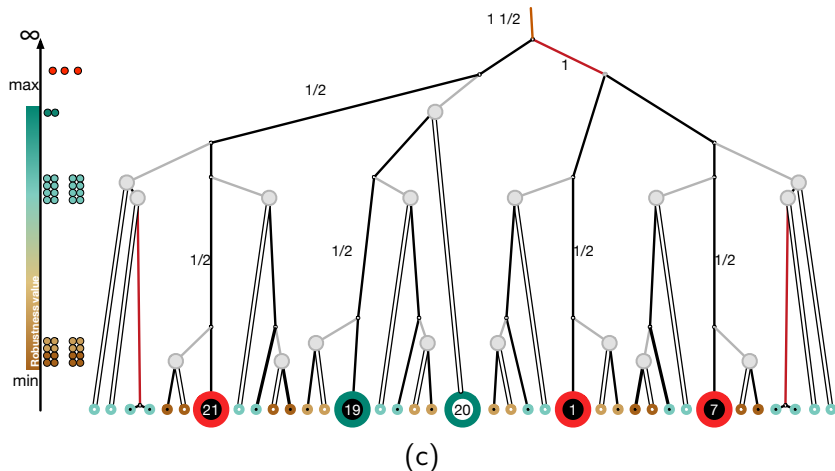


(a)



(b)

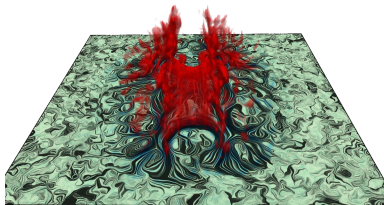
# Stress tensor field: two point load



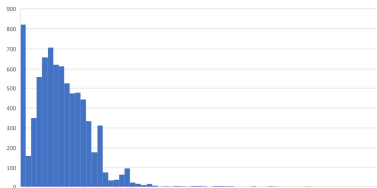
**Figure:** Visualization of the robustness of a slice of a stress tensor field. (c) The full merge tree.



# Diffusion tensor imaging



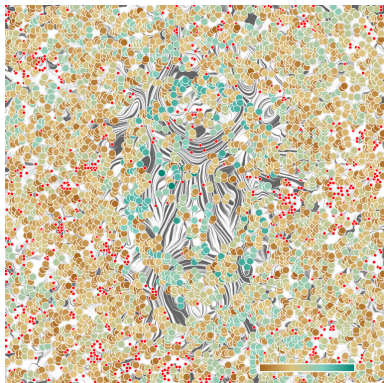
(a)



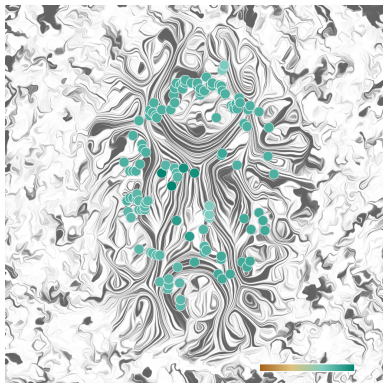
(b)

**Figure:** A slice of a diffusion tensor imaging data set. (a) 3D context visualization using volume rendering of the anisotropy. (b) The histogram of the robustness values of degenerate points.

# Diffusion tensor imaging



(c)



(d)

**Figure:** (c) All degenerate cells are color-coded according to their robustness values. (d) Most robust degenerate cells are highlighted in turquoise.

Thank you for a wonderful and energetic semester!

- Anderson, K. L., Anderson, J. S., Palande, S., and Wang, B. (2018). Topological data analysis of functional mri connectivity in time and space domains. In *Connectomics Neuroimaging (Lecture Notes in Computer Science, Proceedings of International Workshop on Connectomics in NeuroImaging)*, accepted, volume 11083. Springer.
- de Silva, V., Morozov, D., and Vejdemo-Johansson, M. (2009). Persistent cohomology and circular coordinates. *Proceedings 25th Annual Symposium on Computational Geometry*, pages 227–236.
- Jankowai, J., Wang, B., and Hotz, I. (2018). Robust extraction and simplification of 2D tensor field topology. Under review.
- Wang, B. and Hotz, I. (2017). Robustness for 2D symmetric tensor field topology. In Schultz, T., Özarslan, E., and Hotz, I., editors, *Modeling, Analysis, and Visualization of Anisotropy*, pages 3–27. Springer International Publishing.